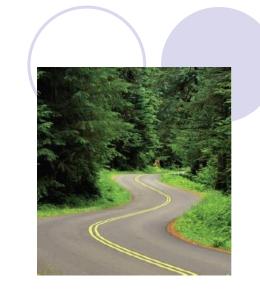
On Optimal Partitioning and Scheduling of DNNs in Mobile Edge/Cloud Offloading

Jie Wu

Dept. of Computer and Information Sciences
Temple University, USA

Roadmap

- 1. On Problem Solving
- 2. Edge/Cloud Computing + AI
- 3. Optimal Scheduling
- 4. Optimal Partition and Scheduling
- 5. Conclusions and Future Work



1. On Problem Solving

How to Solve It (Poyla, 1945)

If you can't solve a problem, then there is an easier problem you can solve: find it.



Is Computing An Experimental Science? (Milner, 1986)

A theory can only emerge through protracted exposure to application.

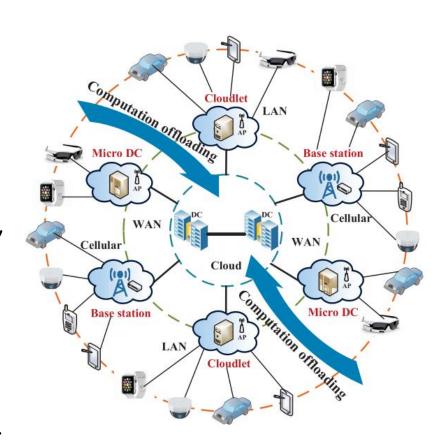
Ideas and applications developed side-by-side



Edge-computing + ML algorithms: traditional solutions

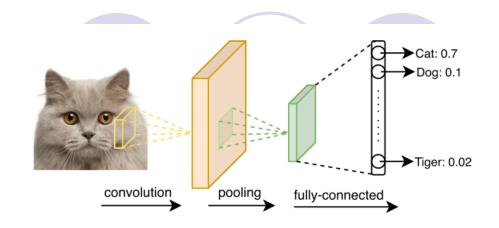
2. Edge/Cloud Computing + AI

- Edge/Cloud Computing
 - Application-driven: AR/VR, video analytics using IoTs
 - Better QoE: mobile/edge device
 - Key indicators: latency, accuracy, energy, and privacy
 - Latency-sensitive
 - How to bring rich computation resources to mobile users?
 - How IoTs contribute to the ML training and inference?

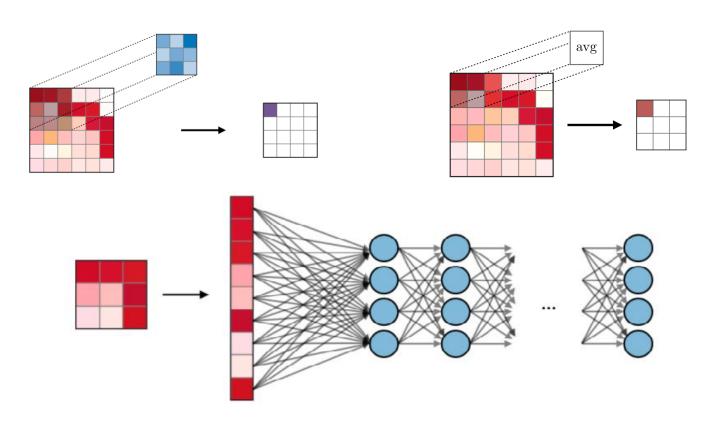


50 billion IoTs: connected intelligence

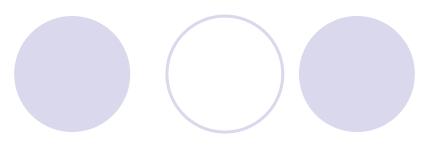
Convolution NNs



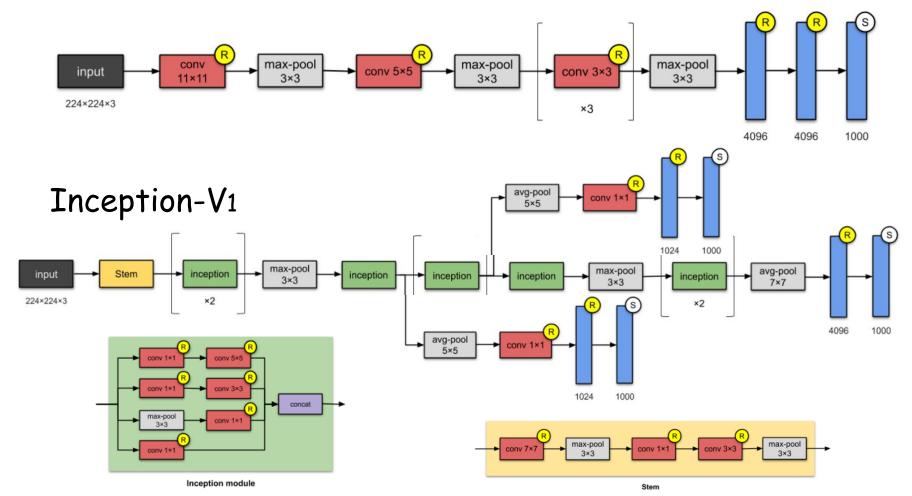
- CNNs (image classification)
- convolution (filtering), pooling (max/avg), fully-connected (neurons)



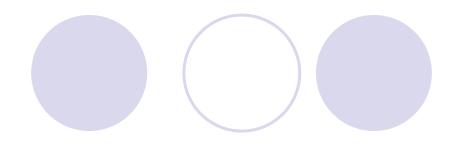
Sample CNNs



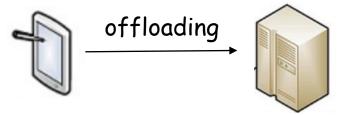
AlexNet (Red: CONV, Gray: POOL, Blue: FC)





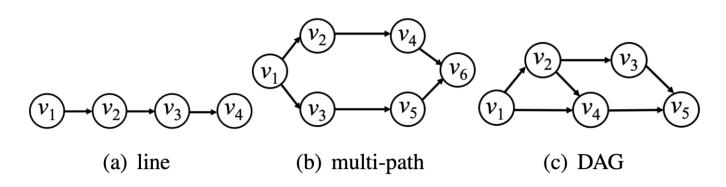


- Three-stage collaborative computation offloading
 - Local computation: processing on local devices
 - Communication: transmitting intermediate DNN layers' outputs
 - Remote computation: completing the remote processing in cloud
- Three models
 - On-device optimization
 - Cloud-only offloading
 - Mixed-mode offloading



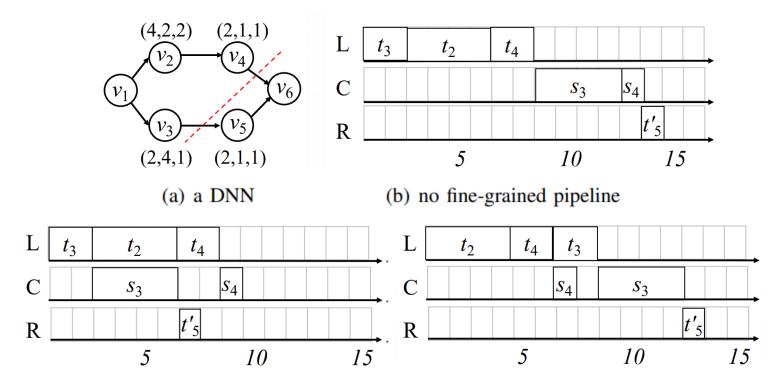
DNN Inferencing

- Deep Neural Networks (DNNs)
 - Technologies: GPU (graphic) and TPU (tensor)
- AI applications
 - Computer vision: AlexNet, VGG-16, Inception, GoogLeNet
 Siamese, Multi-Stream, and RandWire
 - Natural language processing: ChatGPT, GPT-4
- Graph models of DNNs



Offloading Samples

- Given a partition (i.e., cut)
 - Coarse-grained pipeline: local, communication, and remote
 - Fine-grained pipeline: path-based (rather than phase-based)



3. Optimal Scheduling

- DNN Computation Offloading Optimization (DCOO)
 - DCOO: minimum makespan for a given partition (i.e., cut)
- Cases of DNN
 - Line-structure: trivial
 - Multi-path: hard
 - DAG: hard

Theorem 1: DCOO is NP-hard for a multi-path DNN.

Proof: Reduce 3-machine flow-shop to DCOO.

Extended Johnson Algorithm (EJA)

Path p(i) in three stages

 $P_1(i), P_2(i), P_3(i)$

Linear solution (EJA)

- Dividing paths into H and L
- E.g., H = {1}, L = {3, 4, 2}

Algorithm 1 Extended Johnson Algorithm (EJA)

1:
$$H \leftarrow L \leftarrow \phi$$

2: **for**
$$i=1$$
 to m **do**

3: **if**
$$p_1(i) + p_2(i) \le p_2(i) + p_3(i)$$
 then

4:
$$H = H \cup p(i)$$

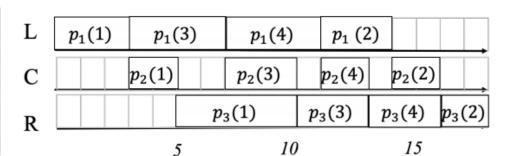
6:
$$L = L \cup p(i)$$

7: Sort H increasingly based on $p_1(i) + p_2(i)$

8: Sort L decreasingly based on $p_2(i) + p_3(i)$

9: Concatenate H and L to obtain σ

Path	$p_1(i)$	$p_2(i)$	$p_3(i)$
i = 1	3	2	5
i = 2	3	2	2
i = 3	4	3	3
i = 4	4	2	3



Optimality

Theorem 2*: If stage 2 is dominated by either stage 1 or 3, $\max\{\min p_1(i), \min p_3(i)\} \ge \max p_2(i)$, EJA is optimal.

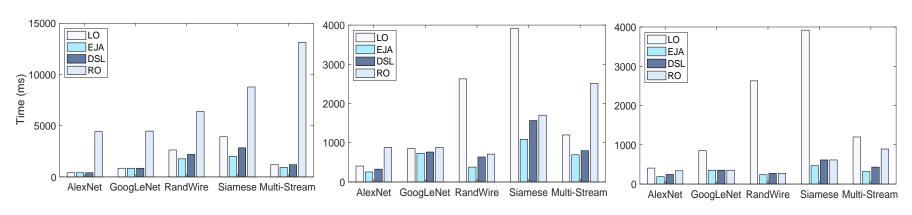
If Theorem 2 fails, EJA still achieves an approximation ratio of 5/3.

Path	$p_1(i)$	$p_2(i)$	$p_3(i)$
i = 1	3	2	5
i = 2	3	2	2
i = 3	4	3	3
i = 4	4	2	3

. .

Simulation

- Local and Cloud
 - Local: Raspberry Pi (and Nexus 4), Cloud: Amazon EC2
 - PyTorch: open-source ML framework
- Algorithms
 - LO: local only, EJA: Extended Johnson's Algorithm,
 DSL: coarse-grained pipeline, RO: remote only



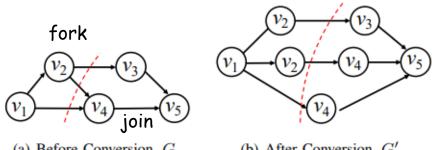
3G (1.1 Mbps)

4G (5.85 Mbps)

WiFi (18.88 Mbps)

Extensions: DAG

- General structure: DAG
 - Conversion to multi-path
 - Replicated nodes at join and fork
- Heuristic solution
 - Scheduling: EJA on multi-path
 - Execution: Replicated node executed once (the first time)



(a) Before Conversion, G

(b) After Conversion, G'

Multiple DNNs Offloading

Internet of Vehicles: smart city

- Autonomous driving systems: perception is a key
- Multiple cameras/sensors: multiple (identical) DNNs
- V2X: V (vehicle); X for I (infrastructure), N (network), or P (pedestrian)



4. Optimal Partition and Scheduling

- Multiple line-structure DNNs
 - AlexNet and VGG-16
 - Video analytics and AR/VR
- Optimal partition and scheduling
 - Brute force: O(kⁿ)
 n: # of copies, k: # of layers
- Existence of a better solution?
 - Exploring special application properties

Johnson Algorithm (JA)

- Closer look at the optimality for EJA
 - o $\max\{\min p_1(i), \min p_3(i)\} \ge \max p_2(i)$
- However, $p_3(i) \approx 0$, reduced to 2-stage pipeline

Algorithm 2 Johnson Algorithm (JA) 1: $H \leftarrow L \leftarrow \phi$ 2: for i = 1 to m do 3: if $p_1(i) \leq p_2(i)$ then

4: $H = H \cup p(i)$

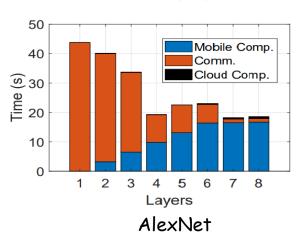
5: else

 $6: L = L \cup p(i)$

7: Sort H increasingly based on $p_1(i)$

8: Sort L decreasingly based on $p_2(i)$

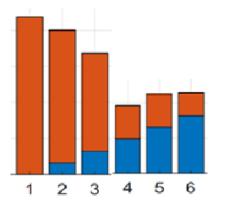
9: Concatenate H and L to obtain σ

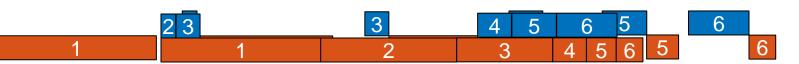


Johnson, Optimal Two- and Three-Stage Production Schedules With Set-up Time Included, Naval Research Logistics Quarter, 1954.

JA in Illustration

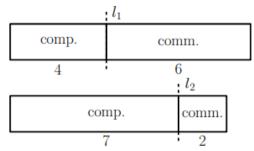
- Optimality is guaranteed: JA on 2-stage pipeline
- First six layers of AlexNet
 - One copy for each partition: 6 copies
 - H = {1, 2, 3}, increasing order of blue(H: comm.-dominate)
 - L = {4, 5, 6}, decreasing order of red(L: comp.-dominate)



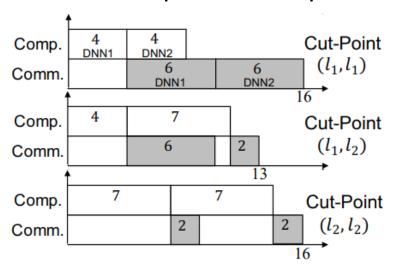


Multiple Line-Structure Example

Two copies of line-structure DNN

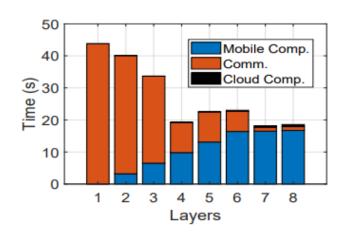


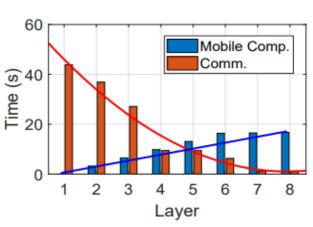
- Three possible partitions and scheduling
 - Gaps in the first and last pairs of comp. and comm.



Special Application Property

- Line-structure
 - Computation time: linear increasing (convex) function
 - Communication time: monotonic decreasing convex function
- Computation vs. communication
 - 🔾 Data size: 2 12 MB
 - Speed (uplink): 2-5 Mpbs (4G) and 6-54 Mpbs (WiFi)





Optimization Approximation

- Two functions
 - Comp. and comm. are convex: one increasing, one decreasing

Theorem 3: A uniform partition of n line DNNs at the intersection will guarantee an approximation of $1 + \frac{1}{n}$.

Proof: convex optimization

- Intersection has the min {max {comp., comm.}}
- Strong duality, then KKT condition, the uniform partition at the intersection has the min max { Σ comp., Σ comm. }
- 1/n is caused by the gaps in the first and last pairs

Duan and Wu, Joint Optimization of DNN Partition and Scheduling for Mobile Cloud Computing, *Proc. of ICPP*, 2021.

Optimization



Pair-wise "merge" and replaced by the middle-point

$$\frac{f(x) + f(x')}{2} \ge f(\frac{x + x'}{2})$$

 The height of the intersection ≤ any max {comp., comm.} of a partition

$$x^*$$

Two gaps, first pair in comm and last pair in comp.: when $n \to \infty, 1 + 1/n$ approaches 1

Insight

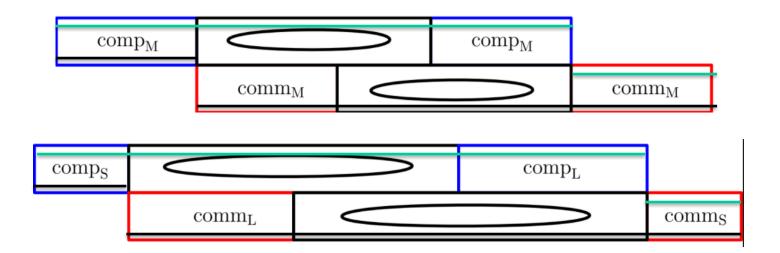
- Comp. (blue) and comm. (orange)
 - o max {blue sum, orange sum}

However, there is a delay gap

Sufficient Condition

- A set of given partitions
 - Left/right most partition: comm, and comp, / comm, and comp,
- Intersection partition: comm_m and comp_m

Theorem 4: The uniform partition beats the given set if $3\text{comp}_m < \text{comp}_s + \text{comp}_l + \text{comm}_s$ and $3\text{comm}_m < \text{comp}_s + \text{comm}_l + \text{comm}_s$



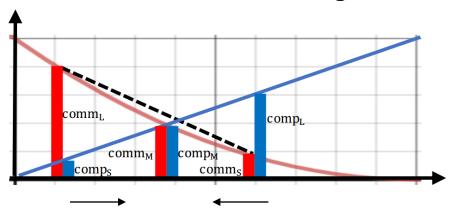
Extended Sufficient Condition

Theorem 5: The uniform partition beats the given set if

(k+2) comp_m < k-prefix.comp_s + k-postfix.comp_l + k-postfix.comm_s

(k+2) comm_m < k-prefix.comp_s + k-prefix.comm_l + k-postfix.comm_s

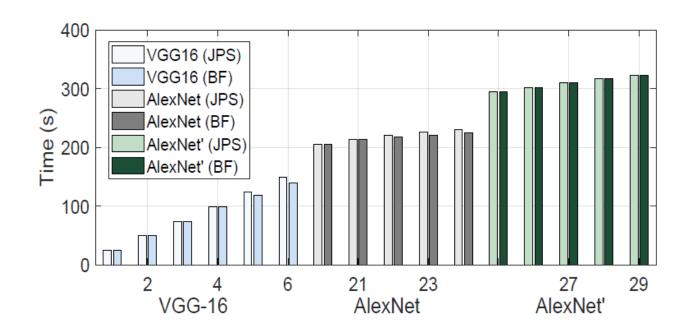
K-prefix/k-postfix: summation of k left/right most partition



Duan and Wu, Optimizing Job Offloading Schedule for Collaborative DNN inference, to appear in *IEEE Transactions on Mobile Computing*, 2023.

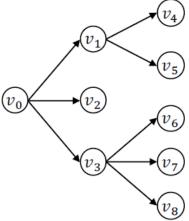
Simulation

- Partition methods
 - Joint Partition and Scheduling: JPS, Brute Force: BF
- Application
 - VGG-16, AlexNet, and AlexNet' (curve fitting) with n = 1, ..., 29



Extension: Tree-structure DNNs

- Merge schedules of subtrees bottom-up (cut at leaves, i.e., $p_3(i) = 0$)
 - Multiway merge of child lists
 - One node at a time, based on Johnson's rule
 - Aggregate computation
 - Root of a subtree and its first child



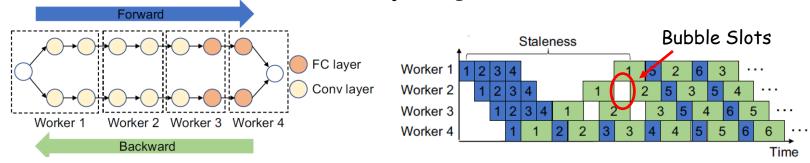
Theorem 6: The schedule generated by the recursive merging approach is optimal for tree-structure DAGs.

Duan and Wu, Computation Offloading Scheduling for Deep Neural Network Inference in Mobile Computing, *Proc. of ACM/IEEE IWQoS*, 2021.

Extension: Inference/Training

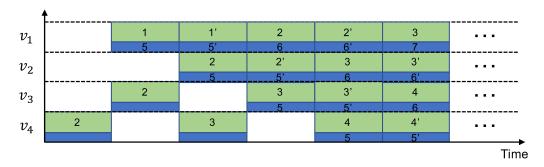
Inference forward pass/training backward pass

Reduce resource idle time by adjusting the ratio of resources



Aligning Pipeline with Resource Allocation

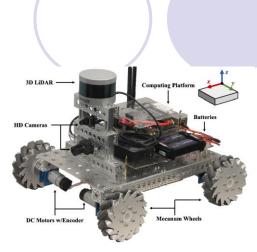
• Combine forward/backward passes (insert 1' after 1 to fill up space)



Duan and Wu, Optimizing Resource Allocation in Pipeline Parallelism for Distributed DNN Training, *Proc. of the IEEE ICPADS*, 2022

An On-going Project

- Extension to DNN training
 - Data compression
- Testbed implementation
 - Visual detection & tracking
- Field test
 - KUSARA at Kettering University



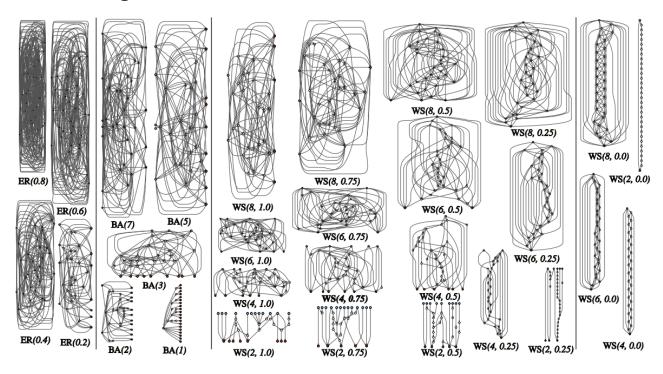


NSF CNS Medium: Cooperative AI Inference in Vehicular Edge Networks for Advanced Driver-Assistance Systems (PI, 2021-2024) (with Stony Brook, Rowan, and Kettering)

Some Reflections

Back to the past: interconnection networks

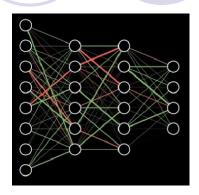
- Randomly wired NNs (random graphs): neuroscience
- Erdos-Renyi (ER): random, Barabasi-Albert (BA): preferential
- Watts-Strogatz (WS): small-world



Xie et al, Exploring Randomly Wired Neural Networks for Image Recognition, *Proc. of ICCV*, 2019.

5. Conclusions and Future Work

- Offloading as a service
 - Mobile Cloud Computing (MCC)
 - ODNN: Single-path, multi-path, and DAG



- Joint partition and scheduling
 - Johnson's rule and its extensions on pipelines
 - Unique properties of comp. and comm. of DNNs
- Future work
 - Optimal partition and scheduling of DAG
 - Pipeline of transfer learning with freeze stage
 - Dynamic nature of offloading speed

Questions



Collaborators: Yubin Duan (Facebook) Ning Wang (Rowan U.)